# **Database Partitioning**

| Database Partitioning              |   |  |
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#### Intro

Consider partitioning as a **fundamental** part of system design even if the system initially only contains a single partition.

A partition is a division of a logical database or its constituent elements into distinct independent parts.

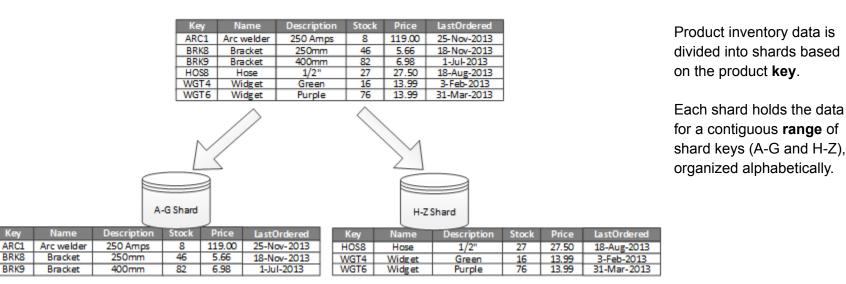
Partitioning is normally done for manageability, performance, availability and load balancing.

It is popular in distributed database management systems.

### Horizontal partitioning / Sharding

Horizontal partitioning / Sharding involves putting different rows into different tables.

Below is an example of horizontal partitioning.



The sharding key must ensure that data is partitioned to spread the workload as evenly as possible across the shards.

The shards don't have to be the same size. It's more important to balance the number of requests.

In practice, we will need to repartition: split large shards, coalesce small shards into larger partitions, or change the schema.

These operations are **time consuming**, and might require taking one or more shards **offline** while they are performed. Replication comes into play.

## Partitioning criteria

|                          | Definition  | Explanation  |
|--------------------------|---|--|
| Range partitioning       | A partition is selected based on column values falling <b>within</b> a given <b>range</b> .       | e.g. A partition for all rows where the column<br>Age has a value between 0 and 18.  |
| List partitioning        | A partition is selected based on column values <b>matching one of</b> a set of discrete values.   | e.g. A partition for all rows where the column<br>City is either Palo Alto, Mountain View,<br>Sunnyvale or San Jose, etc for Santa Clara<br>County.                          |
| Round-robin partitioning | With n partitions, the ith row in insertion order is assigned to the partition <b>(i mod n)</b> . | Pros: sequential access to a relation to be done<br>in parallel<br>Cons: direct access to individual row, based on<br>a predicate, requires accessing the entire<br>relation |
| Hash partitioning        | Applies <b>a hash function</b> to some attribute that yields the partition number                 |  |
| Composite partitioning   | Combinations of the above.  | e.g. Consistent hashing could be considered a composite of hash and list partitioning where the hash reduces the key space to a size that can be listed.                     |

## Vertical partitioning

Vertical partitioning involves putting different **columns** into different tables.

Below is an example of vertical partitioning.

| Key  | Name       | Description | Stock | Price  | LastOrdered |
|------|------------|-------------|-------|--------|-------------|
| ARC1 | Arc welder | 250 Amps    | 8     | 119.00 | 25-Nov-2013 |
| BRK8 | Bracket    | 250mm       | 46    | 5.66   | 18-Nov-2013 |
| BRK9 | Bracket    | 400mm       | 82    | 6.98   | 1-Jul-2013  |
| HOS8 | Hose       | 1/2"        | 27    | 27.50  | 18-Aug-2013 |
| WGT4 | Widget     | Green       | 16    | 13.99  | 3-Feb-2013  |
| WGT6 | Widget     | Purple      | 76    | 13.99  | 31-Mar-2013 |



ARC1 5.66 BRK8 Bracket 250mm 6.98 27.50 400mm BRK9 Bracket 1/2" HOS8 Hose 13.99 WGT4 Widget Green 13.99 WGT6 Widget Purple

Key

| Key  | Stock | LastOrdered |
|------|-------|-------------|
| ARC1 | 8     | 25-Nov-2013 |
| BRK8 | 46    | 18-Nov-2013 |
| BRK9 | 82    | 1-Jul-2013  |
| HOS8 | 27    | 18-Aug-2013 |
| WGT4 | 16    | 3-Feb-2013  |
| WGT6 | 76    | 31-Mar-2013 |

One partition holds Name, Description, Price, which are slow-moving and accessed more frequently.

The other holds Stock, LastOrdered, which are dynamic, commonly used together, and sensitive.

#### Case Study: Apache Cassandra

Cassandra is a horizontally-partitioned row store.

In Cassandra, each node owns a range of tokens.

When data is inserted into the cluster, the first step is to apply a **hash** function to the **partition key**.

The output token is used to determine what node (based on the token range) will get the data.

| OUNTRY | CITY      | POPULATION |             | COUNTRY | СІТҮ      | POPULA  |
|--------|-----------|------------|-------------|---------|-----------|---------|
| AU     | Sydney    | 4.900.000  | Partitioner | 59      | Sydney    | 4.900.  |
| СА     | Toronto   | 6.200.000  |             | 12      | Toronto   | 6.200.0 |
| СА     | Montreal  | 4.200.000  |             | 12      | Montreal  | 4.200.0 |
| DE     | Berlin    | 3.350.000  |             | 45      | Berlin    | 3.350.0 |
| DE     | Nuremberg | 500.000    |             | 45      | Nuremberg | 500.0   |

Cassandra uses **consistent hashing**: it maps every node to one or more **tokens on** a continuous **hash ring**, and defines ownership by hashing a key onto the ring and then "walking" the ring in one direction.

The main difference of consistent hashing to naive data hashing is that when the **number of nodes** to hash into **changes**, consistent hashing only has to **move a small fraction of the keys**.

### Case Study: Apache Spark

Spark is a distributed computing engine. Its low-level data abstraction is a **resilient** distributed dataset (RDD). Resilient because RDDs are **immutable** (can't be modified once created) and **fault tolerant**.

RDD is a distributed collection of objects, which are stored in partitions.

Hash partitioning, range partitioning and custom spark partitioning are supported. Operations such as groupByKey, reduceByKey and sort automatically result in a hash or range partitioned RDD.

Spark creates one task per partition, if there are too many partitions, most of the time goes into creating, scheduling, and managing the tasks then executing. If there are too few partitions, cluster resources may not be fully utilized due to less parallelism.

Relations between partitions of RDDs are categorized as below:

| Narrow dependency  | Wide dependencies  |  |  |
|--|--|--|--|
| each partition of the parent RDD is used by at most one partition of the child RDD   | each partition of the parent RDD is used by multiple partitions of the child RDD |  |  |
| pipelined execution on one node  | shuffled across nodes  |  |  |
| Narrow Transformations<br>1 to 1<br>$\rightarrow$ $\square$<br>$\rightarrow$ $\square$<br>$\rightarrow$ $\square$<br>$\rightarrow$ $\square$ | Wide Transformations (shuffles)<br>1 to N  |  |  |

We may repartition wisely to

## \* References

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